## o Changes

2009/11/04 Corrected for ConTEXt (thanks to Wolfgang Schuster):
Now there's a third party file, $\mathrm{t}-\mathrm{X}_{\exists} \mathrm{Search}$. tex, so that $\mathrm{X}_{\exists}$ Search can be properly loaded with \usemodule[ $X_{\exists}$ Search].
The clash between ConTEXt's \unexpanded macro and $X_{G} T_{E} X^{\prime} s$ (actually $\varepsilon-T_{E} X^{\prime} s$ ) \unexpanded primitive has been fixed.
2009/10/24 Initial version

## 1 Introductory remarks

1. This set of macros requires the $X_{G T} T_{E} X$ engine.
2. This set of macros is totally experimental.
3. This set of macros is written with plain $\mathrm{X}_{-1 \mathrm{~T}} \mathrm{X} \mathrm{X}$, and so it should be compatible with all formats, at least if they implement such basic macros as \newcount or \newif, which is the case at least for $\mathrm{EAT}_{\mathrm{E}} \mathrm{X}$ and $\mathrm{ConT} \mathrm{T}_{\mathrm{E}} \mathrm{Xt}$.
4. As a consequence of the preceding remark, I've used in the examples of this documentation control sequences that don't exist in any format (as far as I know) but whose meaning is transparent enough, like \blue or \italics, which typeset blue and italics. They are not part of $\mathrm{X}_{\exists}$ Search.
5. This set of macros tweaks $\mathrm{X}_{\mathrm{G}} \mathrm{T}_{\mathrm{E}} \mathrm{X}^{\prime} \mathrm{s}$ character class mechanism badly. This mechanism was not designed to do what it does here. Anyway, since it is used mainly for nonalphabetical writing systems, there's little chance of clashing with $X_{\exists} S_{\text {Search. I }}$ I have tried to make $\mathrm{X}_{\exists}$ Search compatible with François Charette's polyglossia for language with special punctuation pattern, like French. I have not tried to patch babel German shorthands in polyglossia, simply because I was not able to make them work
6. $\mathrm{X}_{\exists}$ Search is local all the way down, that is, there's not a single global command. So it can be used in a controlled way. ${ }^{1}$
7. To see what $\mathrm{X}_{\exists}$ Search does, see example 1 on the right.
8. To load the package in thrm{ET}_{\mathrm{E}}\mathrm{X}\),say\usepackage\{xesearch\}InConTEXt:\usemodule[xesearch]Inplain$X_{GTE}T_{E}X$:\inputxesearch.styundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefinedundefined

[^0]\SearchList\{color\}\{\csname\#1\endcsname\{\#1\}\}\{blue, red, green\} This is blue and this is red and this is green, but apparently yellow was not defined

This is blue and this is red and this is green, but apparently yellow was not defined.

## Example 1: A Simple Example

## 2 Let＇s search

－\SearchList $\langle$＊！$\rangle\{\langle$ name $\rangle\}\{\langle$ replacement text $\rangle\}\{\langle$ list of words $\rangle\}$
The star and exclamation mark are optional and their relative order does not matter．Stick－ ing to mandatory arguments for the moment，here＇s how this macro works：first，you give a $\langle$ name $\rangle$ to this list，for further reference．Then you specify the $\langle$ replacement text $\rangle$ ，which will be substituted for all of the words in 〈list of words〉（separated by commas）．In this $\langle$ replacement text $\rangle$ ，the substituted word is designed by \＃1，so just think about it as an ar－ gument to a control sequence．If you forget \＃1，the word disappears（until we learn how to use the exclamation mark），as can be seen in example 2.

Note that there＇s still a space between forgotten and the full stop．Where does it come from？Well，it is the space that was between forgotten and something．At the time when $\mathrm{X}_{\exists}$ Search manipulates something，this space has already been read and typeset，so it does not disappear．

But there＇s something much more interesting in this example．As you might have noticed，the first line says：

```
\SearchList{list1}{\italics{#1}}{obvious7y}
```

and in the text to be searched we find＇Obvious $1 y^{\prime}$＇，with an uppercase first letter．Nonethe－ less，it is found and modified according to the replacement text．We thus discover one basic principle of $\mathrm{X}_{\exists}$ Search：it isn＇t case－sensitive by default．Hence the two following lists

$$
\begin{aligned}
& \text { \SearchList }\{1 \text { ist1\} }\{<\text { whatever }>\} \text { \{word }\} \\
& \text { \SearchList }\{1 \text { ist2\} }\{<\text { whatever }>\}\{\text { Word }\}
\end{aligned}
$$

will find exactly the same set of words，namely＇word＇＇Word＇，＇woRd＇，＇WORD＇，etc．How scary．This isn＇t customary in good programming and in $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ in particular．Fortunately， this default setting can be easily changed：the optional star just after \SearchList will make the entire list case－sensitive．Besides，if a list is not case－sensitive，i．e．if it has no star，a star before a word in that list will make the search for that particular word case－ sensitive．${ }^{2}$ This is illustrated in example 3.

In this example we discover another macro，whose meaning is clear：
－\StopList $\{$ llist of lists〉\}
The lists，separated by commas，are turned off．

[^1]\SearchList\｛1ist1\}\{\italics\{\#1\}\}\{obvious7y\}
\SearchList\｛1ist2\}\{\}\{something\}
Obviously，I have forgotten something．
Obviously，I have forgotten ．

## Example 2：Words As Arguments

```
\SearchList{Case insensitive}{\blue{#1}}{woRd}
Word word woRd WORD
\StopList{Case insensitive}
Word word woRd WORD
\SearchList*{Case sensitive}{\red{#1}}{word}
Word word woRd WORD
\StopList{Case sensitive}
Word word woRd WORD
\SearchList{Mixed}{\green{#1}}{word,*Worm}
Word word woRd WORD\par
Worm worm woRm WORM\par
Word word woRd WORD
Worm worm woRm WORM
```

Example 3：Illustrating Case－Sensitivity

Let's turn back to \SearchList again. It can also take an exclamation mark beside the star (the order between the two of them is not important). In this case, the word is not subsituted anymore; i.e. the replacement text will follow the word (still with \#1 standing for it). These concatenating replacements are very dangerous because they are expanded after the search has started again. You see what I mean: if the word you've found does not endure some transformation that'll make it different from itself as far as typesetting is concerned, ooops, here's the loop. WORD expands to WORD $\backslash$ command\{WORD\} to WORD $\backslash$ command\{WORD $\backslash$ command $\{W O R D\}\}$, etc., and there's no way out of it.

So, what's the point? The point is: the reason why those replacements are placed after the no-search area has stopped is because they are meant to host argument-taking commands to act on the rest of the streams. Such commands can't be placed in normal replacement texts without an exclamation mark, because they would stumble upon precisely what starts the search again. So be careful. Either use !-marked searches with non-typesetting macros, for instance to index the word, or make sure that you know exactly the many interactions you might create. The exclamation mark says it all. Example 4 is silly but I hope you can see the point.

Note the space at the beginning of the first and third replacement texts. Concatenating replacement texts (which replace nothing but whatever) stick to their targets. Besides, in the third example, \green would have gobbled the subsequent space.

I hope you have noticed that the Hamlet list contains not a word but a phrase. So you know: $\mathrm{X}_{\exists}$ Search can find phrases. Now we can't avoid going into a little more detail concerning the way $X_{\exists}$ Search works. But before that, let's see one simple macro:

## - \AddToList $\langle\boldsymbol{*}!\rangle\{\langle$ name $\rangle\}\{\langle$ list of words $\rangle\}$

This adds the words to the $\langle$ name〉 list, which of course should already exist. The presence or absence of a star and/or an exclamation mark doesn't depend at all on the original list. You can see that in example 5.

Finally, the words in \SearchList and \AddToList should be made of characters only, but these can be the product of expansion. For instance, if you have \def $\backslash w o r d\{a$ word\}, then you can say $\backslash$ AddToList $\{m y l i s t\}\{\backslash$ word $\}$. If anything else shows up $\mathrm{X}_{\exists}$ Search won't accept the word (and you'll probably have a good deal of errors beforehand).

## 3 What $\mathrm{X}_{\boldsymbol{g}}$ Search looks for and how it finds it

$\mathrm{X}_{\exists}$ Search can see only two things: letters and non-letters. Non-letters it doesn't like because it's then forced to spit the letters it has gathered and form a word, and most times

```
\SearchList*!{Hamlet}%
    { Or Not \StopSearching#1\StartSearching}%
    {To Be}
To Be...
```

To Be Or Not To Be..
$\backslash$ SearchList! \{typo\}\{\red\{!!!\}\}\{tipo\}
There's a tipo here
There's a tipo!!! here.
\SearchList!\{XeTeX\}\{ \green\}\{is\}
This is \XeTeX.\par
This is $X_{G} T_{E} X$.

## Example 4: A Silly One

\SearchList\{Stupid 1ist\}\{\blue\{\#1\}\}\{word\} Word and beep.
\AddToList*\{Stupid 1ist\}\{Beep\}
Or Beep and word and beep.
Word and beep. Or Beep and word and beep.

## Example 5: Adding Words To An Existing List (Another Silly One)

it's not allowed to take it away. (Un)fortunately, $X_{\exists}$ Search is quite short-sighted: it considers letters what you tell it are not non-letters ( $\mathrm{X}_{\exists}$ Search apparently has some formal education in philosophy).

More seriously (and clearly), $\mathrm{X}_{\exists}$ Search forms a word as long as there are letters. As you can see in example 6, macros are expanded and if they yield letters, $\mathrm{X}_{\exists}$ Search can recognize a word. So when does it stop searching? There are two main cases:

1. It encounters a space, or any primitive control sequence. The former case is quite natural: you want spaces to delimit words (including \skips and associates). But the latter is less obvious: as soon as $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ does something that is not typesetting letters, $\mathrm{X}_{\exists}$ Search gives up. And this includes something as seemingly innocuous as a $\backslash r e l a x$, as you can see in example 7. That's the reason why, for instance, $X_{\exists}$ Search will never find TeX in \TeX: the definition contains many operations that aren't strictly speaking putting letters in the stream. Fortunately, the bulk of a manuscript is made of letters and spaces, and one seldom inserts \relaxes in the middle of words.
2. $\mathrm{X}_{\exists}$ Search encounters a character that you've declared as a non-letter, that is a word boundary. This leads us to the following macro:

- \MakeBoundary\{(characters)\}
- \UndoBoundary\{〈characters〉\}

The characters should be simply put one after the other, as in for instance
$\backslash$ MakeBoundary\{,;:!\}
\UndoBoundary\{?()<br>{<br>}\}

The basic set of such characters is as follows ${ }^{3}$

$$
\text { ., ;:!?-`'()[]\{\} }
$$

Now, if $\mathrm{X}_{\exists}$ Search encounters a character that you've made into a boundary, it will stop forming a word and evaluate what it has gathered. Conversely, such characters cannot appear in the list of words in \SearchList; they wouldn't be found anyway. This is illustrated in example 8.

There is one big difference between those two cases. Characters defined as boundaries are not only word boundaries but also phrase boundaries. If $X_{\exists} S e a r c h$ smells a possible phrase, spaces and primitive commands won't stop it, whereas boundary characters will.

[^2]\SearchList\{Wi11 it find me?\}\{\blue\{\#1\}\}\{word\}
$\backslash d e f \backslash r d\{r d\}$
Here is a wo\rd.
Here is a word.

Example 6: Macros Can't Hide Letters
\SearchList\{This time I'm prepared\}\{\blue\{\#1\}\}\{word\} \def $\backslash r d\{\backslash r e 1 a x$ rd
Here is a wo\rd.
Here is a word.

## Example 7: But Primitive Can

\MakeBoundary\{/\}<br>\SearchList\{separated\}\{\ddag\#1\ddag\} \{waka, jawaka\}<br>Waka/Jawaka<br>$\ddagger W a k a \ddagger / \ddagger J a w a k a \ddagger$<br>\UndoBoundary\{/\}<br>\SearchList\{united\}\{\ddag\#1\ddag\}\{waka/jawaka\}<br>Waka/Jawaka<br>$\ddagger$ Waka/Jawaka $\ddagger$

Example 8: Where Words Start And Stop

You can see that in example 9. This example also illustrates one fact and one sad truth. The fact is that words aren't searched for inside phrases; so the first two you's were not turned to italics, since they belonged to you are what you is. The third one, one the other hand, was recognized since you are neither good nor bad was missed because of the intervenig comma.

The sad truth is that the \kern disappeared. This is one shortcoming of $\mathrm{X}_{\exists}$ Search: primitives disappear when they're in the middle of a possible phrase, even if that phrase is not recognized in the end. By 'possible phrase' I mean a string of words that form the beginning of a phrase that you want identified, e.g. the kern in

```
\SearchList{H(a)unting primitives}{<whatever>}%
    {xesearch feeds on kerns}
xesearch feeds on\kern1cm skips
```

will disappear, even though no string matches in the end. Hopefully such commands are rather rare in the bulk of a document. If some are unavoidable - and for other uses too there exists a pair of commands, whose function I probably don't need to explain (except that $\backslash$ StartSearching doesn't need to be issued at the beginning of your document, it is there by default):

- \StartSearching
- \StopSearching


## 4 (A very blunt form of) regular expressions

Words are cool, and phrases too. But life doesn't always reach their level of achievement. Sometimes you don't know what you want. Something beginning with a 'B', why not? or maybe something that ends in 'et'? Then look at example 10.

There are several things to see in this example. First, $\mathrm{X}_{\exists}$ Search has entered the \italics command and imposed its will. ${ }^{4}$ Next, affixes ${ }^{5}$ are also sensitive to case-sensitivity, so to speak, since beside was not identified (*B? being case-sensitive), whereas PET was found (?et not being case-sensitive). Note that a word matches an affix search if it is at least as

[^3]```
\SearchList{word}{\italics{#1}}{you}
\SearchList{phrases}{\red{#1}}
                                    {you are what you is,
                            you are neither good nor bad}
You are what\kern1cm % What a kern!
you is but you are neither good, nor bad.
You are what you is but you are neither good, nor bad.
```


## Example 9: Phrases And Words

## \SearchList\{Affixes\}\{\red\{\#1\}\}\{*B?,?et,?ET\}

A \italics\{Black Page\} in B, actually some kind of duet for Terry Bozzio and Chad Wackerman, lay on the drumset beside the PET facility.
A Black Page in B, actually some kind of duet for Terry Bozzio and Chad Wackerman, lay on the drumset beside the PET facility.

Example 10: Prefixes And Suffixes
long as the specified part of the affix. Thus, $B$ matches B?. So the question mark means 'from zero to any number of additional letters,' and not 'at least one additional letter.' Phrases can take only suffixes, and they affect the last word only. So

$$
\text { \SearchList\{list\}\{<whatever>\}\{some interesting wor?\} }
$$

will find some interesting world, some interesting words, but not some interesting word thesaurus. An affix mark anywhere else will have no effect.

Marking the unspecified part of a word with ? is the only possibility for the question mark to enter a \SearchList, and obviously it doesn't stand for itself. So, unless of course you undo it as a string boundary, ? can appear only at the beginning or the end of a word. ${ }^{6}$ In any other place, it will be recognized as a boundary that has no right to be there and you'll be blamed. This means that infixes don't exist in $X_{\exists}$ Search, i.e. you can't say B?et to search for bullet, for instance. Also, you can't say ?ull? to match bullet. One affix at a time.

Finally, don't try to use a joker, i.e.
\SearchList\{1ist\}\{<whatever>\}\{?\}
as an attempt to match all words. This won't work. ${ }^{7}$

## 5 Search order(s)

Now we shall see what happens when a word is matched by several searches. There are three different cases:

1. A word is matched by two or more strictly identical searches, e.g.:
```
\SearchList{list1}{<whatever>}{word}
\SearchList{1ist2}{<whatever else>}{word}
... word ...
```

2. A word is matched by two or more prefixes or two or more suffixes identical in casesensitivity, e.g.:
\SearchList\{1ist1\}\{<whatever>\}\{*wor?\}
[^4]If you want to match all words
\SearchList\{1ist\}\{<whatever>\}\{a?,b?, ...,z?\}
should do. Ok, now you've read it, you might have the impression that the title of this section verges on dishonesty. You might be right.
\SearchList\{1ist2\}\{<whatever e1se>\}\{*wo?\}
... word ...
3. A word is matched by two or more different searches, e.g.:

```
SearchList{list1}{<whatever>}{*wor?}
\SearchList{1ist2}{<whatever else>}{word}
\SearchList{1ist3}{<anything>}{?ord}
... word ...
```


### 5.1 Strictly identical searches

In this case, the word will execute all the replacement texts. Their interactions depend on the way they are defined: the replacement texts that are defined without an exclamation mark take as arguments the replacement texts that are defined just before them and will themselves become arguments to subsequent replacement texts. See example 11

If the replacement texts are defined with and exclamation mark, they are simply concatenated, and most importantly, their argument is the word itself alone, not the accumulation of previous remplacement texts. See example 12. Of course, if a word is matched by both kinds of replacement texts, the same rules apply, as in example 13, where you can also be entertained by some not-very-fun-but-you-can-hopefully-see-the-point-again fiddling with !-marked macros. If you want to know what those three \expandafters are doing here, see section 6 .

### 5.2 Affixes with identical characteristics

When a word is found by two or more affixes of the same kind (i.e. only prefixes or only suffixes) and with the same case-sensitivity, then you decide. $X_{\exists}$ Search provides the following commands:

- \SortByLength $\langle *\rangle\{\langle\boldsymbol{p P s S}\rangle\}$
- \DoNotSort $\{\langle p$ PsS $\rangle\}$
- \SearchAll\{〈pPsS〉\}
- \SearchOnlyOne\{ $\{$ pPsS $\rangle\}$
$p, P, s$ and $S$ are shorthands for (respectively) 'case-insensitive prefix', 'case-sensitive prefix', 'case-insensitive suffix' and 'case-sensitive suffix'. They refer to the type of affix to modify and those commands can take one or several of them, e.g. \SearchAll\{pSP\}. By
\SearchList $\{7$ ist1\}\{\blue\{\#1\}\}\{blue word\}
\SearchList\{list2\}\{\dag\#1\dag\}\{blue word\}
\SearchList\{1ist3\}\{\ddag\#1\ddag\}\{blue word\}
This blue word wears earrings and is equivalent
to \ddag\dag\blue\{term\}\dag\ddag.
This $\ddagger+b l u e$ word $\dagger \ddagger$ wears earrings and is equivalent to $\ddagger \dagger t e r m \dagger \ddagger$.


## Example 11: Nested Replacement Texts

> \SearchList!\{1ist1\}\{+\}\{wor?\}
> \SearchList!\{1ist2\}\{\dag\}\{wor?\}
> \SearchList!\{1ist3\}\{\ddag\}\{wor?\}
> This word is a freight train.

This word $+\dagger \ddagger$ is a freight train.

## Example 12: Concatenation <br> (Yet Another Silly Example)

\SearchList\{1ist1\}\{\green\{\#1\}\}\{*?ORD\}
\SearchList\{1ist2\}\{\ddag\#1\ddag\}\{*?ORD\}
\def $\backslash w h i s p e r \# 1\{\backslash i t a l i c s\{(\# 1)\}\}$
$\backslash$ def $\backslash i n g r e e n\{i n$ green\}
\SearchList!\{1ist3\}
$\{\backslash$ expandafter $\backslash$ expandafter $\backslash$ expandafter $\backslash$ whisper \}
\{*?ORD\}
\SearchList!\{1ist4\}\{\ingreen\}\{*?ORD\}
This WORD must be upset.
This $\ddagger$ WORD $\ddagger$ (in green) must be upset

## Example 13: Everything Together

(This Is Mind-Blowing)
default, affixes follow the same rules as full words: each replacement text will take the replacement text defined just before as argument. But you can also create an order between them: with \SortByLength, longer affixes match words before shorter ones, and their replacement texts are thus more deeply nested; adding a star to \SortByLength reverses the order: shorter affixes before longer ones. \DoNotSort resets to default, i.e. replacement texts follow the order in which they were defined. See example 14
$\backslash$ SearchAll and $\backslash$ SearchOnlyOne sets what should happen when a word is matched by an affix: shall the search stop, or shall $\mathrm{X}_{\exists}$ Search continue to investigate whether other affixes might fit too? By default, all affixes are tested, but you might want a different behavior. Thus \SearchOnlyOne\{PS\} will make case-sensitive prefixes and suffixes search only once (and thus the order defined just before becomes extremely important) while \SearchAll\{PS\} will return to default, as illustrated in example 15.

### 5.3 Different searches

Finally, we have to see what $X_{\exists}$ Search should do when several searches match a word. Once again, you decide, thanks to the following command:

- \SearchOrder\{〈order and inhibitions〉\}

You know what p, P, s and S mean; f and F mean 'case-insensitive full word' and 'casesensitive full word.' In the macro above, $\langle$ order and inhibitions $\rangle$ is a list of one or more sequences like f ! ps; (with the semi-colon as part of the expression) in which the red part is optional and which means: if a word matches a full-word case-insensitive search, then $\mathrm{X}_{\exists}$ Search will not test case-insensitive prefixes and suffixes on this word. Such declarations are put one after the other, and this defines the search order. For instance, the default order for $\mathrm{X}_{\exists} \mathrm{Search}^{\text {is: }}$

```
SearchOrder{
    F!fPpSs;
    f!PpSs;
    P!pSs;
    p!Ss;
    S!s;
s;
}
```

and it simply means that full words should be searched for before prefixes, and prefixes before suffixes, with case-sensitive search first in each case, and that any successful search

```
\SearchList{Three letters}{\ddag#1\ddag}{*adv?}
SearchList{Two letters}{\red{#1}}{*ad?}
\SearchList{Four letters}{\dag#1\dag}{*adve?}
\SortByLength{P} adverb
\SortByLength*{P} adverb
\DoNotSort{P} adverb
\ddaggertadverb}\ddagger\ddagger t\ddaggeradverb\ddagger\dagger †\ddaggeradverb\ddagger
```

Example 14: This Is Fascinating
\SearchList\{just a list\}\{\b1ue\{\#1\}\}\{b1?,*bo?\}
\SearchList\{just another list\}\{\bold\{\#1\}\}\{blu?,*bol?\}
\Search0n1y0ne\{P\} B7ue and bold and
\SortByLength\{P\} bold and blue.
Blue and bold and bold and blue.

Example 15: This Guy Sure Ain’t No David Foster Wallace
inhibits any subsequent test. You can have as many sequences as you wish. If $X_{G} T \mathrm{~T} X$ goes crazy and never terminates, then you've probably forgotten a semi-colon (I do it very frequently). See example 16 for an illustration.

Remember that e.g. word? will find 'word' as a prefix, not as a full word, so that 'word' will not be found if you say for instance \SearchList\{list\}\{<whatever>\}\{word?\} and $\backslash$ SearchOrdef $\{\mathrm{f} ;\}$. Finally, although something like $\backslash$ Search $0 r d e r\{f ;\}$ is perfectly okay to search for case-insensitive full words only, $\backslash$ SearchOrder $\{;\}$ will only make $X_{G T E} \mathrm{~T}_{\mathrm{E}} \mathrm{X}$ crazy; \StopSearching is simpler.

## 6 Some TEXnical matters

This section is not vital to the comprehension of $X_{\exists}$ Search, but it may be useful.

- \PrefixFound
- \SuffixFound
- \AffixFound

When a word is found thanks to an affix search, the prefix or suffix used is stored in the relevant macros. If there are several matching affixes, the last prefix and the last suffix win in their respective categories, and between them the same rule apply for $\backslash A f f i x F o u n d$. These macros are available as long as the search has not started again, i.e. they're fully available in normal replacement texts, but in !-marked definitions they're erased as soon as a letter is typeset, so they can be used only at the very beginning. The rest of the time they are empty.

The affix itself respects the case in which it was declared if it is case-sensitive, but it is in lowercase otherwise, however it was fed to \SearchList. See example 17.

- \PatchOutput
- \NormalOutput

By default, $\mathrm{X}_{\exists}$ Search doesn't patch the output routine so footers and headers are searched. This can be done by these two commands. \PatchOutput should of course be issued after any modification to the output routine. \NormalOutput restores the value of the output routine at work when \PatchOutput was executed.

```
\SearchList{word}{\green{#1}}{*Word}
\SearchList{prefix}{\frame{#1}}{wor?}
SearchList{suffix}{\reverse{#1}}{?ord}
\SearchOrder{F;p;s;}
This Word is we11-matched.
\SearchOrder{F!p;p;S;}
This Word is not so wel1-matched anymore.
\SearchOrder{f;}
This Word is not matched at all.
This broW is well-matched.
This Word is not so well-matched anymore.
```

This Word is not matched at all.

## Example 16: Search Order

SearchList\{A case-sensitive suffix\}\{Suf\blue\SuffixFound\}\{*?FiX\} Suffix.

SufFiX.
SearchList\{A case-insensitive affix\}\{\blue\AffixFound fix\}\{Pre?\} PREfix.
prefix

Example 17: Finding Affixes

- \PatchTracing
- \NormalTracing

If you want to give a look at your log file with some tracing on, you will find hundreds if not thousands of totally uninformative lines. That's $\mathrm{X}_{\exists}$ Search recursively discovering new letters and testing words. With $\backslash$ PatchTracing, $X_{\exists}$ Search will try to keep quiet during those painful moments, i.e. \tracingcommands and \tracingmacros will be turned to zero. It can't exactly be totally silent, so just know that all its words begin with xs@. \NormalTracing lets $\mathrm{X}_{\mathrm{G}}$ Search express itself again.

Now just consider example 18. When $\mathrm{X}_{\exists}$ Search reads the input, it introduces itself to all the letters it doesn't know. Most importantly, it writes down some information about them, like their catcode. Now, if a letter is met with a given category catcode, that's the way $\mathrm{X}_{\exists}$ Search will remember it, and this will influence how prefixes and suffixes are recognized. More precisely: the identification of a letter (e.g. the first occurence of it in the typestting stream) and its definition as part of an affix should be done under the same category code.

Note that in example 18 I first had to stop the $f z$ list, otherwise the prefix Frank Zap? would not have been recreated. Another solution would have been to create another prefix like Frank Za? or *Frank Zap?.

Finally, here's how replacement texts are processed. Suppose you have:

```
\SearchList{listone}{\italics{#1}}{word}
\SearchList{listtwo}{\blue{#1}}{word}
\SearchList{1istthree}{\bold{#1}}{word}
```

then $X_{\exists}$ Search does something like this:

```
\def\command@1istone#1{\italics{#1}}
\def\command@1isttwo#1{\blue{#1}}
\def\command@1istthree#1{\bold{#1}}
```

and when word is encountered it is turned to

```
\expandafter\command@1istthree\expandafter{%
    \expandafter\command@1isttwo\expandafter{%
        \expandafter\command@1istone\expandafter{\WORD}}}
```

where \WORD contains exactly word; as you can see, this is equivalent to

```
\command@1istthree{\command@1isttwo{\command@1istone{word}}}
```

which you won't have failed to notice is not equivalent to

```
\bold{\blue{\italics{word}}}
```

```
\catcode`\Z=12
Here's a Z.
\catcode`\Z=11
```

\SearchList\{fz\}\{\italics\{\#1\}\}\{Frank Zap?\}
Look, here comes Frank Zappa!
\StopList\{fz\}
\catcode`\Z=12
\SearchList\{true fz\}\{\italics\{\#1\}\}\{Frank Zap?\}
One more time for the world.
Here comes Frank Zappa!
Here's a Z.
Look, here comes Frank Zappa!
One more time for the world. Here comes Frank Zappa!
Example 18: The Mysterious Z
although in this example the difference is immaterial. Now, if you really want three expansions with superior precision on one word, you probably don't need $\mathrm{X}_{\exists}$ Search: just use a good old macro instead.

Finally, !-marked replacement texts are simply concatenated, as in:

$$
\begin{aligned}
& \text { \expandafter\command@1istone\expandafter }\{\backslash \text { WORD }\} \\
& \text { \expandafter\command@1istthree\expandafter\{\WORD\} } \\
& \text { \expandafter\command@1isttwo\expandafter\{\WORD\} }
\end{aligned}
$$

Now you can see the reason for the three \expandafter's in example 13

## 7 Examples

$\mathrm{X}_{\exists}$ Search was first designed as the basis for the $\mathrm{X}_{\exists}$ Index package, an automatic indexing package for $\mathrm{X}_{\mathrm{G}} \mathrm{ET}_{\mathrm{E}} \mathrm{X}$. It developped into a stand-alone project, and standing so alone that there are no other application yet. So here are some ideas.

First, this document has the following list:
\SearchList*\{logos\}\{\csname\#1\endcsname\}\{?TeX, ?ConTeXt, xesearch\}
(with \xesearch properly defined beforehand) so throughout this document I was able to type 'xesearch can do this or that' to produce ' $X_{\exists}$ Search can do this or that'. That's not fascinating but it was a test.

Being a linguist I can also directly paste examples from my database and turn on $X_{\exists}$ Search to highlight some words. For instance, suppose you're studying the grammaticalization of, say, going to in English, ${ }^{8}$ and you have many examples. Then you just create a command like \startexample, or patch an existing command to activate $\mathrm{X}_{\exists}$ Search just for this stretch of text, among other things. For instance:

```
\(\backslash\) SearchList\{goingto\}\{\bold\{\#1\}\}\{going to\}
\def\startexample\{\%
    Here you can modify margins, for instance.
    \StartSearching
    \}
def \(\backslash\) stopexample\{\%
    \StopSearching
```

[^5]```
Here you restore previous values
```

\}

Otherwise you can locally use \StopList if you're searching the rest of the document too. What follows are some sketchy ideas. Concerning syntax highlighting, I won't try to compete with the listings package.

### 7.1 Spelling

Here's a recipe to create an English spellchecker. Take the list of the 40,000 most frequent words of English by Wiktionary: http://en.wiktionary.org/wiki/Wiktionary: Frequency_lists\#English. Use $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ to turn it into a file, say english.dic, whose only content is \csname<word>@dic \endcsname for each word of the list, with <word> in lowercase. What! you exclaim, that creates 40,000 control sequences! True. But $T_{E} X$ distributions can easily do that today. Input english.dic at the beginning of your document. Then set up $X_{\exists}$ Search as follows:

```
SearchList{spe11ing}{%
    \lowercase{\ifcsname#1@dic\endcsname}%
        #1%
    \else
        \red{#1}%
    \i}
    {a?,b?,c?,d?,e?,f?,g?,h?,i?,j?,k?,1?,m?,
    n?,o?,p?,q?,r?,s?,t?,u?,v?,w?,x?,y?,z?}
\SearchOrder{p;}
```

Now, for each word, $\mathrm{X}_{\exists}$ Search checks whether it belongs to the frequency list. If it doesn't, it puts it in red, thus signaling a likely spelling error. It could also issue an error message, or whatever.

Some words will never belong to that list. Then we use a simple macro to add them beforehand:

```
\def\AddWord#1{\lowercase{\csname#1@dic\endcsname}}
```

We could also create more specific macros like \AddRegularVerb which from e.g. change would add change, changes, changed, changing. $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ could also rewrite english.dic on the fly so there'd be no need to respecify those words on every document. And so on and so forth.

Stately, plump Buck Mulligan came from the stairhead, bearing a bowl of lather on which a mirror and a razor lay crossed. A yellow dressinggown, ungirdled, was sustained gently behind him on the mild morning air. He held the bowl aloft and intoned:

- Introibo ad altare Dei.

Halted, he peered down the dark winding stairs and called out coarsely: - Come up, Kinch! Come up, you fearful jesuit!

Solemnly he came forward and mounted the round gunrest. He faced about and blessed gravely thrice the tower, the surrounding land and the awaking mountains. Then, catching sight of Stephen Dedalus, he bent towards him and made rapid crosses in the air, gurgling in his throat and shaking his head. Stephen Dedalus, displeased and sleepy, leaned his arms on the top of the staircase and looked coldly at the shaking gurgling face that blessed him, equine in its length, and at the light untonsured hair, grained and hued like pale oak.
Buck Mulligan peeped an instant under the mirror and then covered the bowl smartly.

- Back to barracks! he said sternly.

He added in a preacher's tone:

- For this, O dearly beloved, is the genuine Christine: body and soul and blood and ouns. Slow music, please. Shut your eyes, gents. One moment. A little trouble about those white corpuscles. Silence, all.

Example 19: The Words In Red Don’t Belong To The Tor 40,000

Using a list like the frequency list is important because we want all forms of a word to appear; i.e. organized word lists have hear and not hears, because there exists either an algorithm or at least the user's brain to derive hears from hear.

### 7.2 Word count

Another simple use of $\mathrm{X}_{\exists}$ Search is counting words in a document. We define a caseinsensitive list with all letters as prefixes, so all words will be matched (we could add numbers too), as we did in the previous example. Supposing we want words like don't to be counted as one word, then we remove the apostrophe from the word boundaries (in case it signals a dialogue, the following space will delimit the word anyway). And we define the search order as case-sensitive prefixes only, because we don't need anything else. The \shownumber macro is clear, I believe. In the first version of the text on the right it is \let to \relax. It's just for fun.

The \advance on \wordcount has to be \global because there might be (hidden) groups in the text, for instance in font-changing commands.

```
\newcount\wordcount
\def\shownumber{%
    \raise.6\baselineskip\hbox toOpt{\hss\tiny\red{\the\wordcount}}
    }
\SearchList!{wordcount}{\global\advance\wordcount1\shownumber{}}
    {a?,b?,c?,d?,e?,f?,g?,h?,i?,j?,k?,1?,m?,
    n?,o?,p?,q?,r?,s?,t?,u?,v?,w?,x?,y?,z?}
\UndoBoundary{'}
\SearchOrder{p;}
```


### 7.3 Syntax highlighting: $\mathrm{TEX}_{\mathrm{E}}$

At first I'd designed a colorful scheme but it was ugly, so here's something much more sober. We simply create an empty list in which we design a macro to add \stringed primitive commands.

```
\SearchList{hilitex}{\bold{#1}}{}
\def\Add#1{%
    \AddToList{hilitex}{#1}%
    }
```

Stately, plump Buck Mulligan came from the stairhead, bearing a bowl of lather on which a mirror and a razor lay crossed. A yellow dressinggown, ungirdled, was sustained gently behind him on the mild morning air. He held the bowl aloft and intoned:

- Introibo ad altare Dei.

Halted, he peered down the dark winding stairs and called out coarsely: - Come up, Kinch! Come up, you fearful jesuit!

Solemnly he came forward and mounted the round gunrest. He faced about and blessed gravely thrice the tower, the surrounding land and the awaking mountains. Then, catching sight of Stephen Dedalus, he bent towards him and made rapid crosses in the air, gurgling in his throat and shaking his head. Stephen Dedalus, displeased and sleepy, leaned his arms on the top of the staircase and looked coldly at the shaking gurgling face that blessed him, equine in its length, and at the light untonsured hair, grained and hued like pale oak.

There are 158 words.
 bowl smartly.

- Back to ${ }^{173}{ }^{174}$ barracks ${ }^{175}$ ! he said ${ }^{173}$ sternly ${ }^{178}$.

He added in a preacher ${ }^{1898}{ }^{189}{ }^{188}$ tone:
 blood and ouns. Slow music, please. Shut your eyes, gents. One moment.


The total number of words is: 218

## Example 20: Counting Words

```
expandafter\Add\expandafter{\string\def}
expandafter\Add\expandafter{\string\expandafter}
expandafter\Add\expandafter{\string\e1se}
\expandafter\Add\expandafter{\string\fi}
\expandafter\Add\expandafter{\string\e1se}
```

We can't do that for prefixes (and we need them if we want e.g. to underline all userdefined \if), because they would be \stringed and thus of category code 12, which example 18 has shown was a trouble. So we design a macro to add words with a backslash added beforehand. And we use it.

```
def\gobble#1{}
def\AddPrefix#1{%
    \AddToList*{hilitex}{\expandafter\gobble\string\\#1?}%
    }
\AddPrefix{new} \AddPrefix{if}
```

We need one last thing. We want $\backslash$ to be recognized as a letter, because it should be put in bold too. But we also want it to be recognized as a string boundary. The only solution is to make it active and let it expand to \relax (a natural string boundary) plus itself in catcode 12 (which is not defined with $\backslash$ MakeBoundary and is thus a letter for $\mathrm{X}_{\exists}$ Search).

```
\catcode`\|=0
\catcode`\\=13
|def\{|relax|string\}
```

If we pack everything into an usual macro to make verbatim text, then we obtain something along the lines of example 21. Don't forget the typewriter font for the real thrill!

The implementation section of this documentation displays a subtler kind of syntax highlighting, viz. \def and associates put the following command in red and index it too, except commands I don't want to see treated as such, like temporary commands. However, the implementation depends on CodeDoc's macros, so I won't show it here, although you can look at the source

### 7.4 Syntax highlighting: HTML

Coloring HTML is rather easy. The most complicated part concerns word boundaries. $\mathrm{X}_{\exists} S$ earch is used to find elements and attributes. Only case-insensitive full words need to be searched for.

```
\def\mycommand#1 {%
    \expandafter\myothercommand#1%
    \ifwhatever
        \newtoks\mytoks
        \mytoks={...}%
    \e7se
        \mytoks={...}%
    \i
    }
```

Example 21: TEX Highlighted

```
\MakeBoundary{<>/=}
\SearchList{elements}{\bold{\violet{#1}}}
    {htm1,meta,head,body,span,p,div,b,h1,img}
\SearchList{attributes}{\bold{#1}}{align,class,style,src}
\SearchOrder{f;}
< and > delimit markup, so we use them to switch \(X_{\exists}\) Search on and off.
```

```
\catcode `\<=13
```

\catcode `\<=13 \catcode`\>=13
\catcode`\>=13 \def<{\bgroup\catcode`\'=13\catcode`\"=13\char`\<\StartSearching{}}
\def<{\bgroup\catcode`\'=13\catcode`\"=13\char`\<\StartSearching{}} \def>{\egroup\char`\>}

```
\def>{\egroup\char`\>}
```

Quoted text should not be searched, because values to attributes are simply put in blue.
Double quotes and single quotes should exclude each other.

```
\catcode`\"=13
newif\ifdbbegin
\def'{%
    \un1ess\ifsgbegin
        \ifdbbegin \egroup \char`\"
        \else \char`\" \bgroup \dbbegintrue \color{blue}\StopSearching
        \i
    fi
    }
catcode`\'=13
newif\ifsgbegin
\def'{%
    \unless\ifdbbegin
        \ifsgbegin \egroup \char`\'
        \else \char`\'\bgroup \sgbegintrue \color{blue}\StopSearching
        \fi
    \i
    }
```

src and href take links as values, usually underlined. So we do just that.

```
\SearchList!{7inks}{\makelink}{src,href}
```

\def $\backslash$ makelink=\#1\{\%

```
\ifx#1'
    \expandafter\makedbqlink
\e1se
    \expandafter\makesgq7ink
\i
}
\def\makedbq7ink#1"{\StopSearching="\underline{#1}"\StartSearching}
\def\makesgqlink#1'{\StopSearching='\underline{#1}'\StartSearching}
```

The \& . . . ; character denotation is often in red.

```
\catcode`\&=13
\def&#1;{%
    \char`\&
    \red{#1;}%
    }
```

Finally we turn off $\mathrm{T}_{\mathrm{E}}$ X's special characters (quotes are made active by < and >), and we make some useful adjustments.

```
\catcode `\"=12
\catcode`\'=12
\catcode`\#=12
\catcode`\_=12
\catcode `\^=12
\catcode`\%=12
\obeylines
\def\par{\leavevmode\endgraf}
\parindentOpt
```

Example 22 shows the bottom of the CTAN page.
<p>
A perhaps less taxing way to express your appreciation
is to make a
<a href ="https://www.tug.org/donate.html\#ctan">donation</a>\  \— small efforts add up </p>

<div id='footer'><hr />
<table width='100\%'>
<tr>
<td align='left'>
<span id='footer_author'>Site sponsor:
<a href='http://www.tug.org'>TeX Users Group</a></span></td> <td>
<span id='footer_middle'>Internet connection provided by
<a href='http://www.smcvt.edu'>St Michael's College</a></span></td> <td align='right'>
<span id='footer_home'>
<a href='/what_is_ctan.html'>What is CTAN?</a></span></td>
</tr>
</table>
</div>
</body>
</html>

## 8 Implementation

### 8.1 First things first

First we look for $\mathrm{X}_{\mathrm{G}} \mathrm{T} \mathrm{X}$.

These will be used to keep a constant punctuation in spite of catcode-changing packages like babel.

We declare $\mathrm{X}_{\exists}$ Search as a package in $\mathrm{ET}_{\mathrm{E}} \mathrm{X}$.
\unexpanded already exists in $\mathrm{ConT}_{\mathrm{E}} \mathrm{Xt}$, and the meaning of the $\varepsilon-T_{E} X$ primitive is taken over by \normalunexpanded, so we have to make the proper adjustment (many thanks to Wolfgang Schuster, who signalled this to me).
\xs@contextmodule is an empty command let to \relax when $X_{\exists}$ Search is loaded with ConTEXt.
Some keywords, indispensable macros, and a bunch of $\backslash$ new things.

```
\ProvidesPackage{!FileName}[!FileDate!space !FileVersion!space Searching documents.]
\else
\def\MessageBreak{^^J}
\def\xs@err#1{%
    \bgroup
    \newlinechar`\^^J%
    \errorcontextlines=0
    \errmessage{xsearch error: #1}%
        \egroup
    }
\fi
\ifcsname xs@contextmodule\endcsname
\let\xs@unexpanded\normalunexpanded
\else
\let\xs@unexpanded\unexpanded
```

8 \fi
$49 \backslash$ def $\backslash x s @ e n d\{\backslash x s @ e n d\}$
\def $\backslash x s @ e m p t y\}$
$\backslash$ def $\backslash x s @ s t a r\{*\}$
\def $\backslash x s @ e x c l a m a t i o n\{!\}$
\def $\backslash x$ s@@uestion\{?\}
\def $\backslash x s @ s t a r e x c l a m\{*!\}$
\def $\backslash x s @ e x c l a m s t a r\{!*\}$
\def $\backslash x s @ w o r d s\{w o r d s\}$
\def $\backslash x s @ p r e f i x e s\{p r e f i x e s\}$
\def\xs@suffixes\{suffixes\}
59 \def $\backslash x s @ g o b b l e \# 1\}$
60 \def\xs@Lowercase\#1\#2\{\lowercase\{\def\#2\{\#1\}\}\}

61 \let $\backslash x s @ r e l a x \backslash r e l a x$
62 \newcount $\backslash x s @ T e m p C o u n t$
63 \newcount $\backslash x s @ C a s e S e n s i t i v e$
64 \newcount $\backslash x s @ T e m p L e n g t h$
65 \newcount $\backslash x s @ L e n g t h$
66 \newbox $\backslash x s @ B o x$

67 \newif \ifxs@Concatenate
68 \newif \ifxs@String
69 \newif $\backslash i f x s @ A f f i x$
\newif \ifxs@Prefix
\newif \ifxs@Suffix
\newif \ifxs@BadWord
\newif \ifxs@Star
\newif \ifxs@Phrase
\newif $\backslash i f x s @ M a t c h$
\newtoks $\backslash x s @ D e f T o k s$
\newtoks\xs@NoReplaceToks

### 8.2 Character classes

Basic classes: natural delimiters (spaces and primitives), left and right delimiters (set by \MakeBoundary) and the normal class, out of which letters and delimiters will be taken.

This is how we make boundaries. Note that if the character has a character class of 8 or 9 , we don't change it. The interchartoks will be modified, however.

78 \chardef $\backslash x s @ N a t D e l=255$
79 \chardef \xs@lrDel=254
8o \chardef \xs@Classes=253
81 \chardef $\backslash x s @ C l a s s l e s s=0$
82 \XeTeXinterchartoks\xs@lrDel\xs@Classless=\{\xs@LearnLetter\}
$8_{3} \backslash$ XeTeXinterchartoks $\backslash x s @ N a t D e l \backslash x s @ C l a s s l e s s=\{\backslash x s @ L e a r n L e t t e r\}$
$8_{4}$ \XeTeXinterchartoks\xs@NatDel\xs@lrDel\{\xs@EndString\}
$8_{5}$ \xs@TempCount $\backslash x s @ C l a s s e s$
86 \def \xs@Delimiters\{\}
87 \def $\backslash x$ s@MakeDel\#1 $\{\%$
88 \ifx\#1\xs@end
\let\xs@next\relax
\else
\let $\backslash x s @ n e x t \backslash x s @ M a k e D e l$
\unless $\backslash i f n u m$ \the $\backslash$ XeTeXcharclass` \#1 \(=7\) \unless\ifnum\the\XeTeXcharclass`\#1=8 $\backslash$ XeTeXcharclass`\#1=\xs@lrDel \expandafter\def\expandafter\xs@Delimiters\expandafter\{\xs@Delimiters\#1\}\% \fi \fi \(\backslash f i \backslash x s @ n e x t\}\) \xs@MakeDel\\{\\}., ;:!?[()]-'`\xs@end

This is the macro that turn a letter into a letter recording itself. It is recursive. Each new letter is assigned a new character class (from 253 downward), then it is made to start the recording process after delimiters, to stop it before, and to add itself to \xs@String in both case or next to another letter. Before natural delimiters, however, if the word recorded up to now is part of a possible phrase, the process is not stopped. The polyglossia patch is needed when e.g. ? is not turned into a $\backslash x s @ 1 r D e\rceil$ but keeps its character class as defined by polyglossia.
oo \def \MakeBoundary\#1\{\%
101 \xs@MakeDel\#1\xs@end
102 \}
103 \def \UndoBoundary\#1\{\%
\xs@UndoBoundary\#1\xs@end \}
\def $\backslash x$ x@UndoBoundary\#1\{\%
\def\xs@temp\{\#1\}\%
\ifx\xs@temp\xs@end
\let\xs@next\relax
$\backslash e l$ se
\ifnum\the\XeTeXcharclass`\#1=\xs@lrDel \def\xs@RemoveFromDelimiters\#\#1\#1\#\#2\xs@end\{\% \def\xs@Delimiters\{\#\#1\#\#2\}\% \}\% \expandafter\xs@RemoveFromDelimiters\xs@Delimiters\xs@end \fi \XeTeXcharclass`\#1=0
\let\xs@next\xs@UndoBoundary
\fi\xs@next
\}
21 \def\xs@Letters\{\}\%
22 \def\xs@CreateLetter\#1\{\%
\ifx\#1\xs@end
\let\xs@next\relax
$\backslash e l s e$
\expandafter $\backslash$ def $\backslash e x p a n d a f t e r \backslash x s @ L e t t e r s \backslash e x p a n d a f t e r\{\backslash x s @ L e t t e r s \# 1\} \%$
$\backslash$ XeTeXcharclass`\#1=\xs@TempCount
\expandafter\def\csname\the\xs@TempCount @xstring@letter\endcsname\{\#1\}\%
\edef $\backslash x s @ P o l y g l o s s i a P a t c h\{\%$
\xs@unexpanded\{\XeTeXinterchartoks\xs@TempCount7\}\{\%
\xs@unexpanded\{\xdef\xs@String\{\xs@String\#1\}\xs@EndString\}\%
\the\XeTeXinterchartoks0 7\}\%
\xs@unexpanded\{\XeTeXinterchartoks\xs@TempCount8\}\{\%
\xs@unexpanded\{\xdef\xs@String\{\xs@String\#1\}\xs@EndString\}\%

```
    \the\XeTeXinterchartoks0 8}%
    \xs@unexpanded{\XeTeXinterchartoks8\xs@TempCount}{%
        \the\XeTeXinterchartoks8 0 \xs@unexpanded{\xs@StartSring}}%
    }%
\xs@PolyglossiaPatch
\XeTeXinterchartoks\xs@TempCount\xs@Classless{%
    \xdef\xs@String{\xs@String#1}%
    \xs@LearnLetter}%
\XeTeXinterchartoks\xs@lrDel\xs@TempCount{%
    \xs@StopTracing
    \xs@StartString
    }%
\XeTeXinterchartoks\xs@NatDel\xs@TempCount{%
\xs@StopTracing
\xs@StartString
}%
\XeTeXinterchartoks\xs@TempCount\xs@lrDel{%
\xdef\xs@String{\xs@String#1}\xs@EndString}%
\XeTeXinterchartoks\xs@TempCount\xs@NatDel{%
\xdef\xs@String{\xs@String#1}%
\ifcsname\xs@String @xs@phrases@cs\endcsname
\eTeXinterchartokenstate0
\xdef\xs@Stack{%
\xs@String\noexpand\xs@end\xs@unexpanded\expandafter{\xs@Stack}%
}%
\edef\xs@String{\xs@unexpanded\expandafter{\xs@String} }%
\XeTeXinterchartokenstate1
\else
\expandafter\xs@Lowercase\expandafter{\xs@String}\xs@1cString
\ifcsname\xs@lcString @xs@phrases@ncs\endcsname
\XeTeXinterchartokenstate0
\xdef\xs@Stack{%
                    \xs@String\noexpand\xs@end\xs@unexpanded\expandafter{\xs@Stack}%
                    }%
\edef\xs@String{\xs@unexpanded\expandafter{\xs@String} }%
```

This is the recursive macro which creates the $\backslash$ XeTeXinterchartoks for the new letter and all existing letter.
$\mathrm{X}_{\exists}$ Search learns a letter when it encounters a character with character class o. Since \xs@CreateLetter is local, and since it is often executed inside the word box (see $\backslash x s @ S t a r t S t r i n g$ ), we record the letters thus created in \xs@PendingLetters and create them for good after the group.

```
            XeTeXinterchartokenstate
        \else
            \expandafter\expandafter\expandafter\xs@EndString
        \i
        \i
        }%
        \xs@TempCount\xs@Classes
        xs@MakeInterCharToks#1%
        \advance\xs@TempCount-1
        \let\xs@next\xs@CreateLetter
    \fi\xs@next
}
\\def\xs@MakeInterCharToks#1{%
\ifnum\xs@TempCount=\XeTeXcharclass`#1
\XeTeXinterchartoks\xs@TempCount\xs@TempCount {\xdef\xs@String{\xs@String#1}}%
\let\xs@next\relax
\else\let\xs@next\relax
    \expandafter\expandafter\expandafter%
        \xs@Xict\csname\the\xs@TempCount @xstring@letter\endcsname%
        \xs@TempCount{\XeTeXcharclass`#1}%
        \xs@Xict#1{\XeTeXcharclass`#1}\xs@TempCount
        \advance\xs@TempCount-1
        \def\xs@next{\xs@MakeInterCharToks#1}%
    \fi\xs@next}
def\xs@Xict#1#2#3{%
    \XeTeXinterchartoks#2#3{\xdef\xs@String{\xs@String#1}}%
}
\def\xs@PendingLetters{}%
198 \def\xs@LearnLetter#1{%
\xs@CreateLetter#1\xs@end
\ifxs@String
    \xdef\xs@PendingLetters{\xs@PendingLetters#1}%
fi
#1}
```


### 8.3 Search lists

First we define whether there's an! or a * or both.

Then, after a basic check on the name of the list, we record it and defined the macros associated with this list as the second argument; these macros are the normal and !-marked ('noreplace') versions (both are created because there might be an $\backslash$ AddToList of a different type). Finally we launch the

204 \def \SearchList\{\%
205 \xs@ChangeCatcodes
\xs@StarOrExclam\xs@Search
207 \}
208 \def\xs@StarOrExclam\#1\#2\#\{\%

## \def\xs@temp\{\#2\}\%

## \ifx\xs@temp\xs@star

\xs@CaseSensitive2
\xs@Concatenatefalse
\else
\ifx\xs@temp\xs@exclamation
\xs@CaseSensitive0
\xs@Concatenatetrue
\else
\ifx\xs@temp\xs@starexclam
\xs@CaseSensitive2
\xs@Concatenatetrue
\else
\ifx\xs@temp\xs@exclamstar
\xs@CaseSensitive2
\xs@Concatenatetrue
\else
\xs@CaseSensitive0
\xs@Concatenatefalse
\fi
\fi
\fi
\fi\#1\%
\}
233 \def\xs@Search\#1\#2\#3\{\%
\ifcsname\#1@xs@searchlist $\backslash e n d c s n a m e$
\xs@err\{\%
'\#1' already exists. \MessageBreak
Use \string\AddToList\{\#1\}\{<words>\} to add words to it\%
word-maker on the list of words. \AddToList is equivalent with some adjustments.

This takes each word one by one and checks and creates a few things.

3\%
\else
\def\xs@ListName\{\#1\}\%
\expandafter\def\csname\xs@ListName @words\endcsname\{\}\%
\expandafter\def \csname \#1@xs@searchlist\endcsname\#\#1\{\#2\}\%
\expandafter\def\csname \#1@xs@searchlist@noreplace\endcsname\#\#1\{\#2\}\%
\expandafter\xs@MakeWord\#3, \xs@end,\%
\xs@RestoreCatcodes
\fi
\}
\def $\backslash$ AddToList $\{\%$
\xs@ChangeCatcodes
\xs@StarOrExclam\xs@AddToList \}
\def\xs@AddToList\#1\#2\{\%
\ifcsname\#1@xs@searchlist\endcsname
\def\xs@ListName\{\#1\}\%
\expandafter\xs@MakeWord\#2, \xs@end, \%
\xs@RestoreCatcodes
$\backslash e l s e$
xs@err\{`\#1' is not a list\}\% \(\backslash f i\) \xs@RestoreCatcodes \} \def\xs@MakeWord\#1, \{\% \def\xs@TempWord\{\#1\}\% \ifx\xs@TempWord\xs@end \let\xs@next\relax \else \ifcsname\ifnum\xs@CaseSensitive=2*\fi\#1@\xs@ListName\endcsname \xs@err\{You have already specified `\ifnum\xs@CaseSensitive=2*\fi\#1'\%
in '\xs@ListName'. \MessageBreak You can't do it twice\}\%
\else
\csname\#1@\xs@ListName\endcsname
\edef\xs@TempWord\{\#1\}\%

For instance, we parse the word, to find prefixes or suffixes or forbidden things, like control sequences. Then we suppress prefixes and suffixes.

Depending on case-sensitivity, we put the word in lowercase or not, and we define a keyword to record the casesensitivity.

Finally, we patch the replacement texts associated with this word or affix.
$\backslash$ chardef $\backslash x s @$ ParseState=0
\xs@BadWordfalse
\xs@Starfalse
Xxs@Prefixfalse
\xs@Suffixfalse
\xs@ParseWord\#1\xs@end
\unless\ifxs@BadWord
\ifxs@Star
\xs@CaseSensitive1
\expandafter\xs@SuppressPrefix\xs@TempWord $\backslash x s @ e n d$
\fi
\ifxs@Prefix
\expandafter\xs@SuppressSuffix\xs@TempWord
\else
\ifxs@Suffix
\expandafter\xs@SuppressPrefix\xs@TempWord\xs@end
\fi
\fi
\def $\backslash x s @ P h r a s e\} \%$
\ifcase\xs@CaseSensitive
\expandafter\xs@Lowercase\expandafter\{\xs@TempWord\}\xs@TempWord \def $\backslash x s @ c s\{n c s\} \%$
\expandafter\xs@CheckSpaces\xs@TempWord $\backslash x s @ e n d$
\or
$\backslash$ def $\backslash x s @ c s\{c s\} \%$
\expandafter\xs@CheckSpaces\xs@TempWord $\backslash x s @ e n d$
\xs@CaseSensitive0
\or
\def $\backslash x s @ c s\{c s\} \%$
\expandafter\xs@CheckSpaces\xs@TempWord $\backslash x s @ e n d$ \fi
\ifxs@Prefix
\xs@MakePrefix
\def\xs@WordType\{prefixes\}\%
\expandafter\xs@PatchDef\csname\xs@ListName @xs@searchlist\endcsname

This is a basic finite state automaton. It starts in state o. A star brings it in state 1 . In both 0 and 1 , if it finds a letter or a ? it goes in state 2. From there, only letters and a ? at the very end of the word are allowed. Boundaries make it crash. The distinction between stage o and stage 1 is needed just in case the user defines the star as a boundary.

```
        \else
            \ifxs@Suffix
                \xs@MakeSuffix
                \def\xs@WordType{suffixes}%
                \expandafter\xs@PatchDef\csname\xs@ListName @xs@searchlist\endcsname
                \else
                    \def\xs@WordType{words}%
                    \expandafter\xs@PatchDef\csname\xs@ListName @xs@searchlist\endcsname
            \i
        \fi
        \fi
    \i
    \let\xs@next\xs@MakeWord
\fi\xs@next
}
323\def\xs@ParseWord#1{%
\def\xs@temp{#1}%
\ifx\xs@temp\xs@end
    \let\xs@next\relax
    \ifxs@Suffix
        \ifnum\xs@ParseState=3
            \xs@err{You can't have a prefix and a suffix in the same word.\MessageBreak
            '\xs@unexpanded\expandafter{\xs@TempWord}' won't be searched}%
            \xs@BadWordtrue
        \fi
    \fi
    \else
        \let\xs@next\xs@ParseWord
    \expandafter\ifcat\noexpand#1\relax
        \xs@BadChar#1{control sequences are forbidden}%
        \else
            \ifcase\xs@ParseState
            \chardef\xs@TempNum=\XeTeXcharclass`#1 %
            \ifx\xs@temp\xs@star
            \xs@Startrue
```

```
    \chardef\xs@ParseState=1
    \let\xs@next\xs@ParseWord
\else
    \ifx\xs@temp\xs@question
        \xs@Suffixtrue
        \chardef\xs@ParseState=2
        \let\xs@next\xs@ParseWord
        \else
            \ifnum\xs@TempNum>\xs@Classes
                \xs@BadChar#1{it's already a string delimiter}%
        \else
            \chardef\xs@ParseState=2
            \ifnum\xs@TempNum=0
                \xs@CreateLetter#1\xs@end
                    \let\xs@next\xs@ParseWord
                \i
        \f
        \fi
    \f
\or
    \chardef \xs@ParseState=2
    \chardef\xs@TempNum=\XeTeXcharclass`#1 %
    \let\xs@next\xs@ParseWord
    \ifx\xs@temp\xs@question
        \xs@Suffixtrue
    \else
        \ifnum\xs@TempNum>\xs@Classes
        \xs@BadChar#1{it's already a string delimiter}%
        \else
            \ifnum\xs@TempNum=0
            \xs@CreateLetter#1\xs@end
            \let\xs@next\xs@ParseWord
        \fi
        \f
```

```
        \or
            \let\xs@next\xs@ParseWord
            \chardef\xs@TempNum=\XeTeXcharclass`#1 %
            \ifx\xs@temp\xs@question
            \xs@Prefixtrue
            \chardef\xs@ParseState=3
        \else
            \ifnum\xs@TempNum>\xs@Classes
                \xs@BadChar#1{it's already a string delimiter}%
            \else
                    \let\xs@next\xs@ParseWord
            \fi
        \fi
        \or
            \xs@BadChar?{it's already a string delimiter}%
        \fi
    \fi
\i\xs@next
}
\def\xs@BadChar#1#2{%
\def\xs@next##1\xs@end{}%
\xs@BadWordtrue
\xs@err{%
    You can't use '\noexpand#1' in `\xs@unexpanded\expandafter{\xs@TempWord}',\MessageBreak
    #2.\MessageBreak
        -\xs@unexpanded\expandafter{\xs@TempWord}' won't be searched
    }%
}
408\def\xs@CheckSpaces#1\xs@end{%
    \xs@@CheckSpaces#1 \xs@end
4 1 0 ~ \}
411 \def\xs@@CheckSpaces#1 #2\xs@end{%
412 \def\xs@temp{#2}%
\fi
379 \%
\(\backslash 0\)
\chardef\xs@TempNum=\XeTeXcharclass`\#1 \%
\ifx\xs@temp\xs@question
\chardef \xs@ParseState=3
\else
ifnum \xs@TempNum>\xs@Classes
\else
\let\xs@next\xs@ParseWord
\fi
\fi
\xs@BadChar?\{it's already a string delimiter\}\%
\fi
\fi
\}
```

In case the word is a phrase, we have to know that, so we check spaces. In case there are some, we record word1, then word1 word2, then word1 word2 word3, etc., as strings that may lead to phrases and should be recognized as such when $\chi_{\exists}$ Search is searching.

This is in case we find something we don't want in the word.

In case the word was recognized as an affix, we add it to the list of affixes beginning (in the case of prefixes) or ending (in the case of suffixes) with a given letter (this is supposed to make $X_{\exists}$ Search faster: when $X_{\exists}$ Search scans a word, it searches e.g. prefixes if and only if there are prefixes with the same initial letter as the word under investigation, and it compares it to those words only). The affix is also added to the lists sorted by length in both orders.

```
\ifx\xs@temp\xs@empty
    \et\xs@next\relax
\else
    \expandafter\xs@MakePhrase\xs@Phrase\xs@end#1\xs@end
    \def\xs@next{\xs@@CheckSpaces#2\xs@end}%
\i\xs@next
}
\def\xs@MakePhrase#1\xs@end#2\xs@end{%
\ifx\xs@Phrase\xs@empty
    \expandafter\def\csname#2@xs@phrases@\xs@cs\endcsname{}%
    \edef\xs@Phrase{#2}%
\else
    \expandafter\def\csname#1 #2@xs@phrases@\xs@cs\endcsname{}%
        \edef\xs@Phrase{#1 #2}%
    \fi
}%
\def\xs@GetFirstLetter#1#2\xs@end{%
\def\xs@FirstLetter{#1}%
}
\def\xs@MakePrefix{%
\expandafter\ifx\csname\xs@TempWord @\xs@cs @xs@prefixes\endcsname\relax
    \expandafter\xs@GetFirstLetter\xs@TempWord\xs@end
    \ifcsname xs@prefixes@\xs@FirstLetter @\xs@cs\endcsname
        \expandafter\edef\csname xs@prefixes@\xs@FirstLetter @\xs@cs\endcsname{%
            \csname xs@prefixes@\xs@FirstLetter @\xs@cs\endcsname\xs@TempWord,}%
        \def\xs@Sign{<}%
        \xs@Insert{\xs@TempWord}{\csname xs@prefixes@\xs@FirstLetter @\xs@cs @longer\endcsname}%
            \def\xs@Sign{>}%
        \xs@Insert{\xs@TempWord}{\csname xs@prefixes@\xs@FirstLetter @\xs@cs @shorter\endcsname}%
        \else
        \expandafter\edef\csname xs@prefixes@\xs@FirstLetter @\xs@cs\endcsname{\xs@TempWord,}%
        \expandafter\edef\csname xs@prefixes@\xs@FirstLetter @\xs@cs @longer\endcsname{\xs@TempWord,}%
        \expandafter\edef\csname xs@prefixes@\xs@FirstLetter @\xs@cs @shorter\endcsname{\xs@TempWord,}%
        \fi
\fi
```

These suppress the ? at the beginning or the end of the word.
Here's how we sort the list: we check each affix, and we insert the affix to be added just before the the first affix that is shorter or longer, depending on the order.

```
}
\def\xs@GetLastLetter#1{%
\ifx#1\xs@end
    \let\xs@next\relax
\else
    \let\xs@next\xs@GetLastLetter
    \def\xs@LastLetter{#1}%
\fi\xs@next
}
\def\xs@MakeSuffix{%
\expandafter\ifx\csname\xs@TempWord @\xs@cs @xs@suffixes\endcsname\relax
    \expandafter\xs@GetLastLetter\xs@TempWord\xs@end
    \ifcsname xs@suffixes@\xs@LastLetter @\xs@cs\endcsname
            \expandafter\edef\csname xs@suffixes@\xs@LastLetter @\xs@cs\endcsname{%
            \csname xs@suffixes@\xs@LastLetter @\xs@cs\endcsname\xs@TempWord,}%
            \def\xs@Sign{<}%
        \xs@Insert{\xs@TempWord}{\csname xs@suffixes@\xs@LastLetter @\xs@cs @longer\endcsname}%
            \def\xs@Sign{>}%
        \xs@Insert{\xs@TempWord}{\csname xs@suffixes@\xs@LastLetter @\xs@cs @shorter\endcsname}%
        \else
        \expandafter\edef\csname xs@suffixes@\xs@LastLetter @\xs@cs\endcsname{\xs@TempWord,}%
        \expandafter\edef\csname xs@suffixes@\xs@LastLetter @\xs@cs @longer\endcsname{\xs@TempWord,}%
        \expandafter\edef\csname xs@suffixes@\xs@LastLetter @\xs@cs @shorter\endcsname{\xs@TempWord,}%
        \fi
    \fi
    }
\def\xs@SuppressPrefix#1#2\xs@end{\def\xs@TempWord{#2}}
\def\xs@SuppressSuffix#1?{\def\xs@TempWord{#1}}
\def\xs@CountLetter#1{%
\ifx#1\xs@end
    \let\xs@next\relax
\else
    \advance\xs@Length1
    \let\xs@next\xs@CountLetter
\fi\xs@next
}
```

\}
\def \xs@SortList\#1, \{\%
\ifx\#1\xs@end
\edef $\backslash x s @ t e m p l i s t\{\backslash x s @ t e m p l i s t \backslash x s @ T e m p A f f i x\} \$,
\let\xs@next\relax
\else
\xs@Length0
\xs@CountLetter\#1\xs@end
\ifnum\xs@Length\xs@Sign\xs@AffixLength
\edef\xs@templist\{\xs@templist\xs@TempAffix, \#1, \}\%
\let $\backslash x s @ n e x t \backslash x s @ E n d L i s t$
$\backslash e l s e$
\edef\xs@templist\{\xs@templist\#1, \}\%
\let\xs@next\xs@SortList
\fi
\fi\xs@next
\}
\def\xs@EndList\#1\xs@end, \{\%
\edef\xs@templist\{\xs@templist\#1\}\%
\}
\def\xs@Insert\#1\#2\{\%
\def\xs@TempAffix\{\#1\}\%
\xs@Length0
\expandafter\xs@CountLetter\#1 \xs@end
\chardef\xs@AffixLength\xs@Length
\def\xs@templist\{\}\%
\expandafter\expandafter\expandafter\xs@SortList\#2\xs@end,
\expandafter $\backslash l e t \# 2 \backslash x s @ t e m p l i s t$ \}
\def \xs@PatchDef\#1\{\%
\expandafter\edef\csname\xs@ListName @words\endcsname\{\%
\csname\xs@ListName @words\endcsname
\xs@TempWord:::\xs@cs:::\xs@WordType:::\ifxs@Concatenate!\fi:: :\% \}\%
\expandafter\ifx\csname\xs@TempWord @\xs@cs @xs@\xs@WordType\endcsname\relax\%

Stopping a list is a delicate process: we have to extract the definition associated with the list from the words where it appears, and it is nested in case it is not !-marked.
\xs@DefToks\{\xs@FinalString\}\%
\else
\xs@DefToks $\backslash$ expandafter $\backslash$ expandafter $\backslash$ expandafter $\{\%$
\csname\xs@TempWord @\xs@cs @xs@\xs@WordType\endcsname\}\%
$\backslash f i$
\expandafter\ifx\csname\xs@TempWord @\xs@cs @xs@\xs@WordType @noreplace\endcsname\relax
\xs@NoReplaceToks\{\}\%
\else
\xs@NoReplaceToks \expandafter \expandafter\expandafter $\{\%$
\csname\xs@TempWord @\xs@cs @xs@\xs@WordType @noreplace\endcsname\}\%
$\backslash f i$
\ifxs@Concatenate
\expandafter\edef\csname\xs@TempWord @\xs@cs @xs@\xs@WordType\endcsname\{\the\xs@DefToks\}\%
\expandafter\edef\csname\xs@TempWord @\xs@cs @xs@\xs@WordType @noreplace\endcsname\{\%
\the\xs@NoReplaceToks
\xs@unexpanded\{\expandafter\#1\expandafter\{\xs@String\}\}\%
3\%
\else
\expandafter\edef\csname\xs@TempWord @\xs@cs @xs@\xs@WordType\endcsname\{\%
\noexpand\expandafter\noexpand\#1\noexpand\expandafter\{\the\xs@DefToks\}\%
\}\%
$\backslash f i$
\}
\def $\backslash$ StopList\{\%
\xs@ChangeCatcodes
\xs@StopList
\}
\def $\backslash x s @ S t o p L i s t \# 1\{\%$
\xs@@StopList\#1, \xs@end,\%
\xs@RestoreCatcodes
\}
\def\xs@@StopList\#1, \{\%
\def\xs@temp\{\#1\}\%
\ifx\xs@temp\xs@end
\let\xs@next\relax
We modify the adequate replacement text: no-replace or normal.
Removing from no-replace is rather easy, since it's nothing more than:
\expandafter\<7ist1-macro>\expandafter\{\xs@String\}

```
\else
```

\else
\ifcsname\#1@xs@searchlist\endcsname
\ifcsname\#1@xs@searchlist\endcsname
\unless\ifcsname\#1@xs@stoppedlist\endcsname
\unless\ifcsname\#1@xs@stoppedlist\endcsname
\csname\#1@xs@stoppedlist\endcsname
\csname\#1@xs@stoppedlist\endcsname
\expandafter\def\expandafter\xs@ToRemove\expandafter{%
\expandafter\def\expandafter\xs@ToRemove\expandafter{%
\csname\#1@xs@searchlist\endcsname
\csname\#1@xs@searchlist\endcsname
}%
}%
\expandafter\expandafter\expandafter%
\expandafter\expandafter\expandafter%
\xs@PatchWords\csname \#1@words\endcsname\xs@end:::::::::::::%
\xs@PatchWords\csname \#1@words\endcsname\xs@end:::::::::::::%
\i
\i
\else
\else
\xs@err{`#1' is not a list}%         \xs@err{`\#1' is not a list}%
\fi
\fi
\let\xs@next\xs@@StopList
\let\xs@next\xs@@StopList
\fi\xs@next
\fi\xs@next
}
}
def\xs@PatchWords\#1:::\#2:::\#3:::\#4:::{%
def\xs@PatchWords\#1:::\#2:::\#3:::\#4:::{%
\def\xs@TempWord{\#1}%
\def\xs@TempWord{\#1}%
\ifx\xs@TempWord\xs@end
\ifx\xs@TempWord\xs@end
\let\xs@next\relax
\let\xs@next\relax
\else
\else
\def\xs@temp{\#4}%
\def\xs@temp{\#4}%
\ifx\xs@temp\xs@exclamation
\ifx\xs@temp\xs@exclamation
\expandafter\expandafter\expandafter%
\expandafter\expandafter\expandafter%
\xs@RemoveFromNoReplace\expandafter\xs@ToRemove\csname\#1@\#2@xs@\#3@noreplace\endcsname
\xs@RemoveFromNoReplace\expandafter\xs@ToRemove\csname\#1@\#2@xs@\#3@noreplace\endcsname
\f
\f
\def\xs@cs{\#2}%
\def\xs@cs{\#2}%
\def\xs@WordType{\#3}%
\def\xs@WordType{\#3}%
\expandafter\xs@RemoveFromDef\csname\#1@\#2@xs@\#3\endcsname
\expandafter\xs@RemoveFromDef\csname\#1@\#2@xs@\#3\endcsname
\let\xs@next\xs@PatchWords
\let\xs@next\xs@PatchWords
\fi\xs@next
\fi\xs@next
}
}
\def\xs@RemoveFromNoReplace\#1\#2{%
\def\xs@RemoveFromNoReplace\#1\#2{%
\def\xs@Erase\#\#1\expandafter\#1\expandafter\#\#2\#\#3\xs@end{%
\def\xs@Erase\#\#1\expandafter\#1\expandafter\#\#2\#\#3\xs@end{%
\def\#2{\#\#1\#\#3}%

```
    \def#2{##1##3}%
```

\expandafter\<list2-macro>\expandafter\{\xs@String \expandafter\<7ist3-macro>\expandafter\{\xs@String\} So we define a macro on the fly to find the definition we want to remove. If there's nothing left, we let this no-replace to $\backslash r e l a x$, so this word might be removed altogether when we evaluate what we find.
Normal replacement texts have the following structure:
\expandafter $\backslash$ <list1-macro> $\backslash$ expandafter
\expandafter\<list2-macro>\expandafter\{

$$
\ldots
$$

\xs@Fina1String

## \}\}

So we scan this recursively and rebuild it piecewise, removing the list that was stopped. If in the end there remains \xs@FinalString only, then there's no replacement text anymore, and if moreover the no-replace part is equal to \relax, then there's nothing left for that word and it shouldn't be tested anymore. So we let the definition associated with this word to \relax or we remove it from affixes.

```
    ifx#2\xs@empty
    \let#2\relax
    \fi
    }%
\expandafter\xs@Erase#2\xs@end
}
def\xs@final{\xs@FinalString}
\def\xs@TempDef{}
\def\xs@RemoveFromDef#1{%
\def\xs@TempDef{}%
\def\xs@Def{\xs@FinalString}%
\unless\ifx#1\xs@final
    \expandafter\xs@Extract#1%
\fi
\let#1\xs@Def
\ifx#1\xs@final
    \expandafter\ifx\csname\expandafter\xs@gobble\string#1@noreplace\endcsname\relax
            \ifx\xs@WordType\xs@words
                \let#1\relax
            \else
                \xs@RemoveFromAffixes
            \fi
        \fi
\fi
}
def\xs@Extract\expandafter#1\expandafter#2{%
\def\xs@temp{#1}%
\unless\ifx\xs@temp\xs@ToRemove
    \edef\xs@TempDef{%
        \noexpand#1,%
        \xs@unexpanded\expandafter{\xs@TempDef}%
        }%
    \fi
\def\xs@temp{#2}%
\ifx\xs@temp\xs@final
```

```
    \def\xs@next{%
    \expandafter\xs@Rebuild\xs@TempDef\xs@end,%
            }%
\else
    \def\xs@next{%
    \xs@Extract#2%
    }%
\fi\xs@next
}
\def\xs@Rebuild#1,{%
\ifx#1\xs@end
    \let\xs@next\relax
    \else
        \let\xs@next\xs@Rebuild
        \edef\xs@Def {%
        \xs@unexpanded{\expandafter#1\expandafter}%
        \noexpand{%
        \xs@unexpanded\expandafter{\xs@Def}%
        \noexpand}%
        3%
    \fi\xs@next
}%
\def\xs@RemoveFromAffixes{%
\ifx\xs@WordType\xs@prefixes
    \expandafter\xs@GetFirstLetter\xs@TempWord\xs@end
    \let\xs@Letter\xs@FirstLetter
\else
    \expandafter\xs@GetLastLetter\xs@TempWord\xs@end
    \let\xs@Letter\xs@LastLetter
\i
\def\xs@templist{}%
\expandafter\expandafter\expandafter%
    \xs@CleanList\csname xs@\xs@WordType @\xs@Letter @\xs@cs\endcsname\xs@end,%
\expandafter\let\csname xs@\xs@WordType @\xs@Letter @\xs@cs\endcsname\xs@templist
\def\xs@templist{}%
```

```
\expandafter\expandafter\expandafter%
    \xs@CleanList\csname xs@\xs@WordType @\xs@Letter @\xs@cs @shorter\endcsname\xs@end,%
\expandafter\let\csname xs@\xs@WordType @\xs@Letter @\xs@cs @shorter\endcsname\xs@templist
\def\xs@templist{}%
\expandafter\expandafter\expandafter%
    \xs@CleanList\csname xs@\xs@WordType @\xs@Letter @\xs@cs @longer\endcsname\xs@end,%
\expandafter\let\csname xs@\xs@WordType @\xs@Letter @\xs@cs @longer\endcsname\xs@templist
\expandafter\let\csname\xs@TempWord @\xs@cs @xs@\xs@WordType\endcsname\relax
66 }
\def\xs@CleanList#1,{%
\def\xs@temp{#1}%
\ifx\xs@temp\xs@end
    \let\xs@next\relax
\else
    \let\xs@next\xs@CleanList
    \unless\ifx\xs@temp\xs@TempWord
        \edef\xs@templist{\xs@templist#1,}%
    \fi
    \i\xs@next
}
```


### 8.4 Testing words

Here comes the big part: collecting words and testing them. When a letter follows a delimiter, we reset some values and start collecting the letters in a box...

681 \xs@Stringtrue
682 \let\xs@StartString\relax
683 \def\xs@String\{\}\%
684 \def $\backslash$ PrefixFound\{\}\%
685 \def $\backslash$ SuffixFound\{ $\} \%$
686 \def $\backslash$ AffixFound\{\}\%
687 \def\xs@Stack\{\}\%
688 \def $\backslash x s @ R e m a i n d e r\} \%$
689 \xs@Phrasefalse
690 \setbox\xs@Box=\hbox\bgroup
...and when a delimiter shows up again, unless we're track ing a phrase, we close the box, create the unknown letters that we've found in it, evaluate the word and finally output the result of this evaluation.

And here are the tests. The F test is for case-sensitive full words and just checks whether there is a definition for this word in this case. If it finds anything, it puts it around the string that already exists, i.e. either the bare word or the word alreay surrounded by replacement texts. Hence The bunch of \expandafters. If there's a no-replace, we also add it to the existing ones. \xs@relax is just a placeholder to add the inhibitions defined with $\backslash$ SearchOrder.

The f does the same thing, except it puts the word in lowercase before hand.
\}
\let\xs@@StartString\xs@StartString
\def $\backslash x s @ E n d S t r i n g\{\%$
\ifxs@String
\egroup
\xs@Stringfalse
\expandafter\xs@CreateLetter\xs@PendingLetters\xs@end
\gdef\xs@PendingLetters\{\}\%
\xs@Evaluate
\xs@Restore
\xs@StartTracing
\expandafter\xs@Remainder
$\backslash f i$
\}
\def $\backslash x s @ @ F @ T e s t\{\%$
\expandafter\unless\expandafter\ifx\csname\xs@String @cs@xs@words\endcsname\relax
\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\% $\backslash d e f \%$
\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\% \xs@FinalString\%
\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\{\% \csname\xs@String @cs@xs@words\endcsname\}\%
\expandafter\unless\expandafter\ifx\csname\xs@String @cs@xs@words@noreplace\endcsname\relax $\backslash e d e f \backslash x s @ N o R e p l a c e\{\%$
\xs@unexpanded $\backslash$ expandafter\{\xs@NoReplace\}\%
\xs@unexpanded $\backslash$ expandafter $\{\backslash$ csname\xs@String @cs@xs@words@noreplace\endcsname\}\%
\}\%
\fi
\xs@Matchtrue
\xs@relax
\xs@relax
\fi
\}
\def \xs@@f@Test\{\%
\expandafter\xs@Lowercase\expandafter\{\xs@String\}\xs@1cString

Tests on prefixes check whether there exists a prefix list beginning with the same letter as the word at stake, and in this case run the \xs@CheckPrefixes test.
\expandafter\unless\expandafter\ifx\csname\xs@lcString @ncs@xs@words\endcsname\relax
 $\backslash d e f \%$
\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\% \xs@FinalString\%
\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter $\backslash e x p a n d a f t e r\{\%$ \csname\xs@lcString @ncs@xs@words\endcsname\}\%
\expandafter\unless\expandafter\ifx\csname\xs@lcString @ncs@xs@words@noreplace\endcsname\relax \edef $\backslash x s @ N o R e p l a c e\{\%$
\xs@unexpanded\expandafter\{\xs@NoReplace\}\%
\xs@unexpanded\expandafter\{\csname\xs@lcString @ncs@xs@words@noreplace\endcsname\}\% \}\%
\fi
\xs@Matchtrue
\xs@relax
\xs@relax

## \fi

\}
def $\backslash x s @ @ p @ T e s t\{\%$
\xs@Affixfalse
\expandafter\xs@GetFirstLetter\xs@lcString\xs@end
\ifcsname xs@prefixes@\xs@FirstLetter @ncs\endcsname
\let\xs@@String\xs@lcString
\def\xs@cs\{ncs\}\%
\let\xs@WhatNext\xs@p@WhatNext
\expandafter\expandafter \expandafter\%
\xs@CheckPrefixes\csname xs@prefixes@\xs@FirstLetter @ncs\p@order\endcsname\xs@end,\% \fi
\ifxs@Affix
\xs@Affixfalse
\xs@Matchtrue
\xs@relax
\xs@relax
\fi
\}

Prefixes are tested one by one by creating a macro on the fly where one delimiter is the prefix. Then we put the word at stake before it and execute the macro, and if there's no first argument, then the word matches the prefix. For instance, if the word is democracy and the prefix is demo then we test \xs@TestPrefix democracydemo
and obviously the first argument is empty, since demo is a delimiter.

```
\def\xs@@P@Test{%
\xs@Affixfalse
\expandafter\xs@GetFirstLetter\xs@String\xs@end
\ifcsname xs@prefixes@\xs@FirstLetter @cs\endcsname
        \let\xs@@String\xs@String
    \def\xs@cs{cs}%
        \let\xs@WhatNext\xs@P@WhatNext
        \expandafter\expandafter\expandafter%
        \xs@CheckPrefixes\csname xs@prefixes@\xs@FirstLetter @cs\P@order\endcsname\xs@end,%
    \fi
    \ifxs@Affix
        xs@Affixfalse
        \xs@Matchtrue
        \xs@relax
        \xs@relax
        \i
    }
\def\xs@CheckPrefixes#1,{%
    \def\xs@temp{#1}%
    \ifx\xs@temp\xs@end
        \let\xs@next\relax
    \else
        \def\xs@TestPrefix##1#1##2\xs@end{%
            \def\xs@temp{##1}%
            \ifx\xs@temp\xs@empty
                \xs@Affixtrue
                \def\PrefixFound{#1}%
                \def\AffixFound{#1}%
            \let\xs@next\xs@WhatNext
            \expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter%
                    \def%
            \expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter%
                \xs@FinalString%
            \expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter{%
                \csname#1@\xs@cs @xs@prefixes\endcsname}%
```

The tests for suffixes work along the same lines as those for prefixes.

```
    \expandafter\unless\expandafter\ifx\csname#1@\xs@cs @xs@prefixes@noreplace\endcsname\relax
        \edef\xs@NoReplace{%
            \xs@unexpanded\expandafter{\xs@NoReplace}%
            \xs@unexpanded\expandafter{\csname#1@\xs@cs @xs@prefixes@noreplace\endcsname}%
            }%
        \i
        \else
            \let\xs@next\xs@CheckPrefixes
        \f
        }%
        \expandafter\xs@TestPrefix\xs@@String#1\xs@end
\i\xs@next
}
```


# \expandafter\unless\expandafter\ifx\csname\#1@\xs@cs @xs@prefixes@noreplace\endcsname\relax 

``` \edef\xs@NoReplace\{\%
\xs@unexpanded\expandafter\{\csname\#1@\xs@cs @xs@prefixes@noreplace\endcsname\}\% \}\%
\fi
\let\xs@next\xs@CheckPrefixes
\(\backslash f i\)
\}\%
fi\xs@next
8 \}
809 \def\xs@@S@Test\{\%
\xs@Affixfalse
\expandafter\xs@GetLastLetter\xs@String\xs@end
\ifcsname xs@suffixes@\xs@LastLetter @cs\endcsname
\let\xs@@String\xs@String
\def\xs@cs\{cs\}\%
\let\xs@WhatNext\xs@S@WhatNext
\expandafter \expandafter\expandafter\%
\xs@CheckSuffixes\csname xs@suffixes@\xs@LastLetter @cs\S@order\endcsname\xs@end,\%
\fi
\ifxs@Affix
\xs@Affixfalse
\xs@Matchtrue
\xs@relax
\xs@relax
\fi
\}
826 \def\xs@@s@Test\{\%
\xs@Affixfalse
\expandafter\xs@GetLastLetter\xs@lcString\xs@end
\ifcsname xs@suffixes@\xs@LastLetter @ncs\endcsname
\let\xs@@String\xs@lcString
```

```
    \def\xs@cs{ncs}%
    \let\xs@WhatNext\xs@s@WhatNext
    \expandafter\expandafter\expandafter%
    xs@CheckSuffixes\csname xs@suffixes@\xs@LastLetter @ncs\s@order\endcsname\xs@end,%
\fi
\ifxs@Affix
    \xs@Affixfalse
    \xs@Matchtrue
    \xs@relax
    \xs@relax
    \fi
}
\def\xs@CheckSuffixes#1,{%
\def\xs@temp{#1}%
\ifx\xs@temp\xs@end
    \let\xs@next\relax
\else
\def\xs@TestSuffix##1#1##2\xs@end{%
\def\xs@@temp{##2}%
\ifx\xs@temp\xs@@temp
\xs@Affixtrue
\def\SuffixFound{#1}%
\def\AffixFound{#1}%
\let\xs@next\xs@WhatNext
\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter%
\def%
\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter%
\xs@FinalString%
\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter\expandafter%
{%
\csname#1@\xs@cs @xs@suffixes\endcsname}%
\expandafter\unless\expandafter\ifx\csname#1@\xs@cs @xs@suffixes@noreplace\endcsname\relax
\edef\xs@NoReplace{%
\xs@unexpanded\expandafter{\xs@NoReplace}%
\xs@unexpanded\expandafter{\csname#1@\xs@cs @xs@suffixes@noreplace\endcsname}%
```

```
866
868
870
83
```

3\%
\fi
\else\%
\let\xs@next\xs@CheckSuffixes
\fi
\}\%
\expandafter\xs@TestSuffix\xs@@String\#1\xs@end
\fi\xs@next
\}

### 8.5 Search order

\SearchOrder actually defines \xs@Evaluate. First it adds inhibitions to the tests, e.g. 'F!f;' adds \7et $\backslash x s @ f @ T e s t \backslash r e 7 a x 876$ to the F test in case it is positive, then it adds the tests them- 877 selves, in the specified order, to \xs@Evaluate.

If the stack is not empty, it means we're dealing with a phrase; so the evaluation is not over in case no test has succeded. We first have to test the phrase minus the last word, then

```
879 \def\xs@SearchOrder#1{%
\def\xs@Order{}%
\xs@@Search0rder#1\xs@end;%
\edef\xs@Evaluate{%
    \xs@unexpanded{%
            \XeTeXinterchartokenstate=0
            \def\xs@NoReplace{}%
            \let\xs@FinalString\xs@String
            \expandafter\xs@Lowercase\expandafter{\xs@String}\xs@lcString
            }%
            \xs@unexpanded\expandafter{%
            \xs@Order
            \ifxs@Match
            \def\xs@next{%
                    \xs@FinalString
            }%
        \else
            \unless\ifx\xs@Stack\xs@empty
            \xs@Phrasetrue
            \expandafter\xs@PopStack\xs@Stack\xs@@end
```

877 878 \}

```
def \SearchOrder{%
\xs@ChangeCatcodes
    xs@SearchOrder
    *x
            ef\xs@next{%
```

the phrase minus the last two words, etc.

If the word was not a phrase, and no test was successful, we simply put the box that contains it back into the stream.

We initialize the tests.

This treats each specification in $\backslash$ SearchOrder and the inhibitions, if any.

```
            \let\xs@next\xs@Evaluate
        \else
            \ifxs@Phrase
                                    \def\xs@Stack{}%
                    \def\xs@next{\xs@String\xs@Restore}%
            \else
                \def\xs@next{\unhbox\xs@Box\xs@Restore}%
            \fi
        \fi
        \i\xs@next
        }%
    }%
\let\xs@f@Test\xs@@f@Test
\let\xs@F@Test\xs@@F@Test
\let\xs@p@Test\xs@@p@Test
\let\xs@P@Test\xs@@P@Test
\let\xs@s@Test\xs@@s@Test
\let\xs@S@Test\xs@@S@Test
\xs@RestoreCatcodes
}
\def \xs@@SearchOrder#1#2;{%
\def\xs@temp{#1#2}%
\ifx#1\xs@end
    \let\xs@next\relax
\else
    \def\xs@Inhibit{}%
    \xs@MakeInhibit#2\xs@end
    \expandafter\expandafter\expandafter\xs@PatchTest\csname xs@@#1@Test\endcsname#1%
    \edef\xs@Order{%
        \xs@unexpanded\expandafter{\xs@Order}%
        \xs@unexpanded\expandafter{\csname xs@#1@Test\endcsname}}%
    \let\xs@next\xs@@SearchOrder
    \i\xs@next
}
933\def\xs@MakeInhibit#1{%
```

The evaluation ends in any case with the restoration of the tests, in case they were inhibited. the remainder is the right part of a discarded phrase. For instance, if $X_{\exists}$ Search searches for page layout it will investigate page properties if it finds it, and the remainder is properties.

```
\def\xs@temp{#1}%
\ifx#1\xs@end
    \let\xs@next\relax
\else
    \let\xs@next\xs@MakeInhibit
    \unless\ifx\xs@temp\xs@exclamation%
        \edef\xs@Inhibit{%
            \xs@unexpanded\expandafter{\xs@Inhibit
            \expandafter\let\csname xs@#1@Test\endcsname\relax}%
        }%
    \fi
\i\xs@next
}
\def\xs@PatchTest#1\xs@relax#2\xs@relax#3#4{%
\expandafter\edef\csname xs@@#4@Test\endcsname{%
    \xs@unexpanded{#1}%
    \xs@unexpanded\expandafter{\expandafter\xs@relax\xs@Inhibit\xs@relax\fi}%
    }%
}
|xs@Restore{%
    \x@Matchfalse
    \let\xs@f@Test\xs@@f@Test
    \let\xs@F@Test\xs@@F@Test
    \let\xs@p@Test\xs@@p@Test
    \let\xs@P@Test\xs@@P@Test
\let\xs@s@Test\xs@@s@Test
\let\xs@S@Test\xs@@S@Test
\let\xs@StartString\xs@@StartString
\edef\xs@Remainder{%
    \xs@unexpanded\expandafter{\xs@NoReplace}%
    \xs@unexpanded\expandafter{\xs@Remainder}%
    }%
\XeTeXinterchartokenstate=1
968 \def \xs@PopWord\#1\xs@end\#2\xs@end\{\%
```

967 \}
eration. The stack itself is built when the beginning of a phrase is found before a natural delimiter.

To search affixes in a given order, we simply define the list to be used in tests to be the one with this order.

```
\def\xs@String{#2}%
\def\xs@@PopWord#2##1\xs@end{%
    \edef\xs@Remainder{##1\xs@unexpanded\expandafter{\xs@Remainder}%
        }%
    3%
\xs@@PopWord#1\xs@end
}
\def\xs@PopStack#1\xs@end#2\xs@@end{%
\def\xs@Stack{#2}%
\expandafter\xs@PopWord\xs@String\xs@end#1\xs@end
}
980 \def\SortByLength#1 {%
\def\xs@temp{#1}%
\ifx\xs@temp\xs@star
    \def\xs@AffixOrder{@shorter}%
    \let\xs@next\xs@SortByLength
    \else
        \def\xs@AffixOrder{@longer}%
        \def\xs@next{\xs@@SortByLength#1\xs@end}%
    \fi
    \xs@next}%
\def\xs@SortByLength#1{%
    \xs@@SortByLength#1\xs@end
}
\def\xs@@SortByLength#1{%
\ifx#1\xs@end
    \let\xs@next\relax
\else
    \expandafter\let\csname #1@order\endcsname\xs@AffixOrder
        \let\xs@next\xs@@SortByLength
    \fi\xs@next
}
1001 \def\DoNotSort{%
1002 \def\xs@AffixOrder{}%
1003 \xs@SortByLength
```

Searching all affixes is done by setting the \xs@WhatNext macro to \xs@<affix>@WhatNext, depending on the text being performed.

Searching only one affix is simply gobbling the remaining ones in case of a successful test.

004 \}
1005 \def $\backslash$ SearchAll\#1 $\{\%$
1006 \xs@SearchAll\#1\xs@end
1007 \}
1008 \def $\backslash x s @$ SearchAll\#1 1 \%
1009 \ifx\#1\xs@end
$1010 \backslash$ let $\backslash x s @ n e x t \backslash r e l a x$
1011 \else\let\xs@next\xs@SearchAll
1012 \if\#1p\%
1013 \let $\backslash x s @ p @ W h a t N e x t \backslash x s @ C h e c k P r e f i x e s$
1014 \else
\else
\if\#1P
\let\xs@P@WhatNext\xs@CheckPrefixes
\else
\if\#1s
\let\xs@s@WhatNext\xs@CheckSuffixes
$\backslash e l s e$
\let $\backslash x s @ S @ W h a t N e x t \backslash x s @ C h e c k S u f f i x e s$
\fi
\fi
$\backslash f i$
\fi\xs@next
\}
1027 \def $\backslash$ SearchOnlyOne\#1\{\%
1028 \xs@SearchOne\#1\xs@end
1029 \}
$1030 \backslash$ def $\backslash x s @ S e a r c h O n e \# 1\{\%$
1031 \ifx\#1 \xs@end
032 \let\xs@next\relax
\else
\let\xs@next\xs@SearchOne
\expandafter\def\csname xs@\#1@WhatNext\endcsname\#\#1\xs@end,\{\}\%
$\backslash f i \backslash x s @ n e x t$
\}

### 8.6 Miscellanea

For the moment, starting and stopping the search is quite brutal.

Patching the output very simple too.

```
1038 \def\StopSearching{%
1039 \let\xs@StartString\relax
1040 }
1041 \def\StartSearching{%
1042\let\xs@StartString\xs@@StartString
1043 }
1044 \let\xs@OldOutput\relax
1045 \def\PatchOutput{%
1046 \ifx\xs@OldOutput\relax
1047 \edef\xs@PatchOutput{%
1048 \noexpand\def\noexpand\xs@OldOutput{%
\the\output
            3%
            \noexpand\output{%
                \noexpand\StopSearching
                \the\output
                    \noexpand\StartSearching
                    }%
                }%
                \expandafter\xs@PatchOutput
    \else
        \xs@err{Output already patched}%
    \i
    }
1062 \def\NormalOutput{%
1063 \ifx\xs@OldOutput\relax
        \xs@err{Output has not been patched}%
    \else
        \expandafter\output\expandafter{%
            \xs@OldOutput
            }%
        \let\xs@OldOutput\relax
    \i
    }
```

As is patching the tracing.
finally we set everything back to normal, set some default values and say goodbye.

1072 \def $\backslash$ PatchTracing\{\%
1073 \def \xs@StopTracing\{\%
1074 \chardef \xs@tracingcommands \tracingcommands
1075 \chardef \xs@tracingmacros\tracingmacros
1076 \tracingcommands0 \tracingmacros0 $\backslash$ relax \}\%
\def\xs@StartTracing\{\%
\tracingcommands\xs@tracingcommands
\tracingmacros\xs@tracingmacros
3\%
1082 \}
1083 \def $\backslash$ NormalTracing\{\%
1084 \let \xs@StopTracing\relax
1085 \let\xs@StartTracing\relax
1086 \}
1087 \NormalTracing
1088 \xs@RestoreCatcodes \catcode`@=12
1089 \SearchOrder\{
1090 F!fPpSs;
1091 f!PpSs;
1092 P!pSs;
1093 p!Ss;
1094 S!s;
1095 s;
1096 \}
1097 \DoNotSort\{pPsS\}
1098 \SearchAll\{pPsS\}
1099 \XeTeXinterchartokenstate1
1100 \endinput

| $1 \% \mathrm{D}$ | \module |
| :--- | ---: |
| $2 \% \mathrm{D}$ | $\mathrm{r} \quad$ file= ! FileName, |
| $3 \% \mathrm{D}$ | version=!FileDate, |
| $4 \% \mathrm{D}$ | title $=$ CONTEXT $\backslash$ User Module, |

as tests for $\mathrm{ConT}_{\mathrm{E}} \mathrm{Xt}$, because there might exist commands with the same names in other formats.)
$6 \%$ D
7 \% D
$8 \%$ D
9 \%D
$10 \% \mathrm{D}$
11
12 \writestatus\{loading\}\{ConTeXt User Module / XeSearch\}
13 \csname xs@contextmodule\endcsname
14 \input xesearch.sty
15 \endinput
subtitle=XeSearch
author=Paul Isambert,
date=\currentdate,
copyright=Paul Isambert,
email=zappathustra@free.fr,
license=LaTeX Project Public License]

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[^0]:    ${ }^{1}$ If your knowledge of $T_{E} X$ is confined to ${ }^{\operatorname{LA}} \mathrm{T}_{\mathrm{E}} \mathrm{X}$, you might not be very familiar with the notion of locality to groups, since in ${ }^{L T} T_{E} \mathrm{X}$ pretty much everything is global by default, whereas in plain $\mathrm{T}_{\mathrm{E}} \mathrm{X}$ the contrary holds. So to make things simple, just remember that if you use $X_{\exists} S e a r c h$ inside a $\mathrm{LA}_{\mathrm{E}} X$ environment, even one you've defined yourself with \newenvironment, nothing will spread outside this environment. (I don't know the situation for ConTEXt, so I won't say anything.)

[^1]:    ${ }^{2}$ However，if $\backslash$ SearchList is suffixed with a star，all words in the list will be case－sensitive．

[^2]:    ${ }^{3}$ That is: full stop, comma, semi-colon, colon, exclamation mark, question mark, dash, inverted comma, apostrophe (i.e. left and right quote), parentheses, brackets, curly braces. This is rather arbitrary, despite some basic sensible assumptions.

[^3]:    ${ }^{4}$ Provided I'm using commands that don't cancel each other, like plain $T_{E} X^{\prime}$ 's $\backslash$ bf and $\backslash$ it
    ${ }^{5}$ I use the word affixes to refer to both prefixes (like B?) and suffixes (like ?et). From a linguistic point of view, prefixes and suffixes (and infixes, actually) are indeed affixes, but from the same point a view, what we're talking about here has nothing to do with prefixes or suffixes, just with bits of words. I hope you don't mind.

[^4]:    ${ }^{6}$ And if a star is present, it should precede the question mark

[^5]:    ${ }^{8}$ If you're a linguist, I apologize for my lack of originality.

